

Randomized Algorithms (RAs)
■ Randomized algorithms are frequently used in many areas of engineering, computer science, physics, finance, optimization,but their appearance in systems and control is mostly limited to Monte Carlo simulations
■ Main objective of this mini-course: Introduction to rigorous study of RAs for uncertain systems and control, with specific applications
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Randomized Algorithms (RAs)

- Combinatorial optimization, computational geometry
- Examples: Data structuring, search trees, graph algorithms, sorting (RQS), ...
- Motion and path planning problems
- Mathematics of finance: Computation of path integrals
- Bioinformatics (string matching problems)

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- Uncertainty has been always a critical issue in control theory and applications
- First methods to deal with uncertainty were based on a stochastic approach
- Optimal control: LQG and Kalman filter
- Since early 80's alternative deterministic approach (worst-case or robust) has been proposed

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Robustness

- Major stepping stone in 1981: Formulation of the \mathcal{H}_{∞} problem by George Zames
- Various "robust" methods to handle uncertainty now exist: Structured singular values, Kharitonov, optimization-based (LMI), *l*-one optimal control, quantitative feedback theory (QFT)

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Robustness

- Late 80's and early 90's: Robust control theory became a well-assessed area
- Successful industrial applications in aerospace, chemical, electrical, mechanical engineering, ...
- However, ...

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Limitations of Robust Control - 1

- Researchers realized some drawbacks of robust control
- Consider uncertainty Δ bounded in a set \mathcal{B} of radius ρ . Largest value of ρ such that the system is stable for all $\Delta \in \mathcal{B}$ is called (worst-case) robustness margin
- Conservatism: Worst case robustness margin may be small
- Discontinuity: Worst case robustness margin may be discontinuous wrt problem data

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Limitations of Robust Control - 2

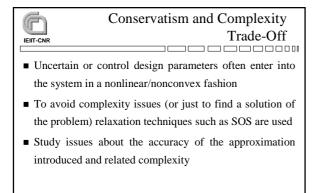
- Computational Complexity: Worst case robustness is often \mathcal{NP} -hard (not solvable in polynomial time unless $\mathcal{P} = \mathcal{NP}$)^[1]
- Various robustness problems are *NP*-hard
 - static output feedback
 - structured singular value
 - stability of interval matrices

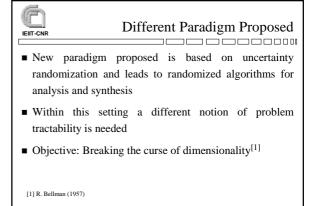
[1] V. Blondel and J.N. Tsitsiklis (2000)

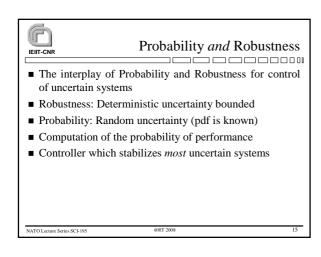
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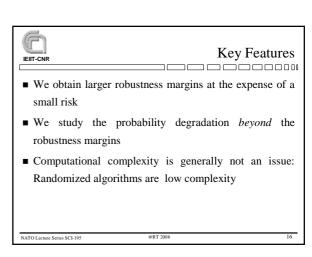
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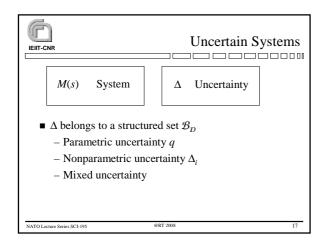
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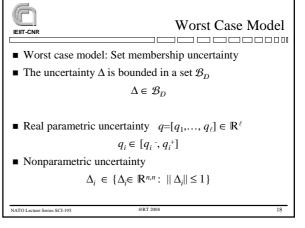


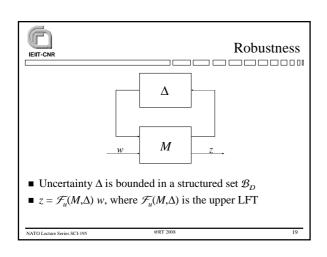


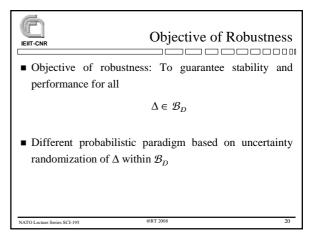


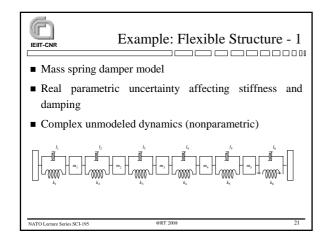


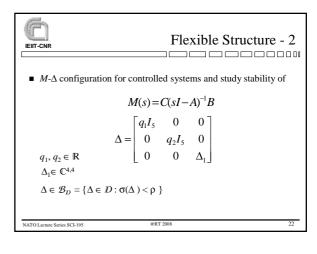


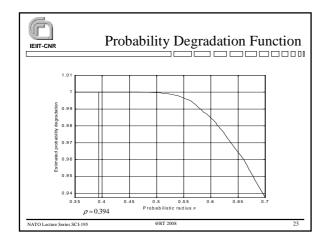


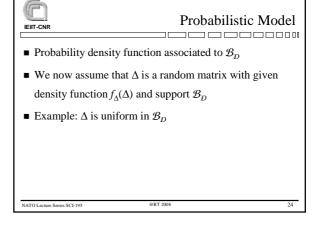














Uniform Density

■ Take $f_{\Delta}(\Delta) = \mathcal{U}[\mathcal{B}_D]$ (uniform density within \mathcal{B}_D)

$$\mathcal{U}[\mathcal{B}_{\!\scriptscriptstyle D}] = \begin{cases} \frac{1}{vol(\mathcal{B}_{\!\scriptscriptstyle D})} & \text{if } \Delta \in \mathcal{B}_{\!\scriptscriptstyle D} \\ 0 & \text{otherwise} \end{cases}$$

■ In this case, for a subset $\mathcal{S} \subseteq \mathcal{B}_D$

$$\Pr\{\Delta \in \mathcal{S}\} = \frac{\int_{\mathcal{S}} d\Delta}{vol(\mathcal{B}_D)} = \frac{vol(\mathcal{S})}{vol(\mathcal{B}_D)}$$

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Performance Function

- In classical robustness we guarantee that a certain performance requirement is attained for all $\Delta \in \mathcal{B}_D$
- This can be stated in terms of a performance function

$$J = J(\Delta)$$

■ Examples: \mathcal{H}_{∞} performance and robust stability

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Example: \mathcal{H}_{∞} Performance - 1

$$J(\Delta) = || \mathcal{F}_{u}(M, \Delta) ||_{\infty}$$

■ For given γ >0, check if

$$J(\Delta) < \gamma$$

for all Δ in \mathcal{B}_D

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Example: $\mathcal{H}_{\scriptscriptstyle\infty}$ Performance - 2

lacktriangle Continuous time SISO systems with real parametric uncertainty q with upper LFT

$$\mathcal{F}_{u}(M,\Delta) = \mathcal{F}_{u}(M,q) =$$

$$\frac{0.5q_1q_2s+10^{-5}q_1}{(10^{-5}+0.05q_2)s^2+\left(0.00102+0.5q_2\right)s+(2\cdot 10^{-5}+0.5q_1^2)}$$

where $q_1 \in \, [0.2,\, 0.6]$ and $q_2 \in \, [10^{\text{-5}}, 3 \cdot 10^{\text{-5}}]$

- Letting $J(q) = \| \mathcal{F}_u(M,q) \|_{\infty}$, we choose $\gamma = 0.003$
- Check if $J(q) < \gamma$ for all q in these intervals

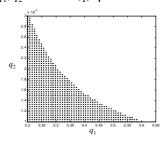
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Example: \mathcal{H}_{∞} Performance - 3

■ The set of q_1 , q_2 for which $J(q) < \gamma$ is shown below



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Example^[1]: Robust Stability - 1

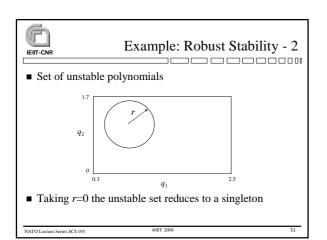
■ Consider the closed loop uncertain polynomial

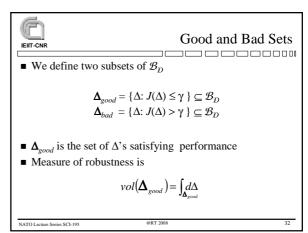
$$p(s,q) = (1+r^2+6q_1+6q_2+2q_1q_2)+(q_1+q_2+3)s+(q_1+q_2+1)s^2+s^3$$
 where $q_1 \in [0.3, 2.5], q_2 \in [0,1.7]$ and r =0.5

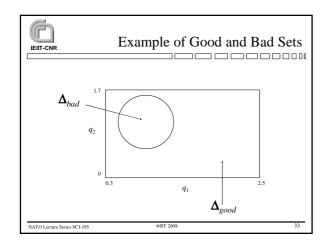
• Check stability for all q in these intervals

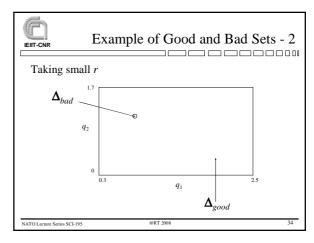
[1] G. Truxal (1961)

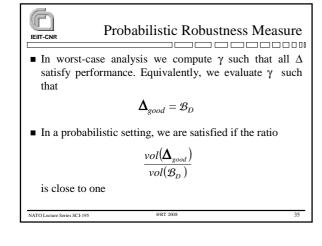
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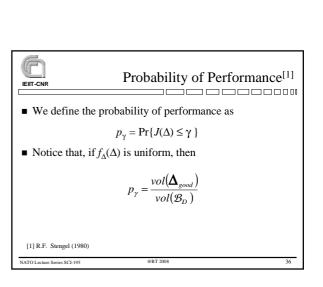














Example: Closed-Form Computation

- For Truxal's example, we compute p_{γ} in closed-form
- For uniform distribution, we have

$$vol(\Delta_{good}) = 3.74 - \pi r^2$$

 $vol(\mathcal{B}_D) = 3.74$

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P1: Performance Verification

■ For given performance level γ , check whether

$$J(\Delta) \le \gamma$$

for all Δ in \mathcal{B}_D

■ Compute the probability of performance p_{γ}

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P2: Worst-Case Performance

■ Find J_{max} such that

$$J_{\max} = \max_{\Delta \in \mathcal{B}_D} J(\Delta)$$

■ Compute the worst case performance (or its probabilistic counterpart)

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Randomized Algorithms for Analysis

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Randomized Algorithm: Definition

- Randomized Algorithm (RA): An algorithm that makes random choices during its execution to produce a result
- Example of a "random choice" is a coin toss

head

or

tails



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Randomized Algorithm: Definition

- Randomized Algorithm (RA): An algorithm that makes random choices during its execution to produce a result
- For hybrid systems, "random choices" could be switching between different states or logical operations
- For uncertain systems, "random choices" require (vector or matrix) random sample generation

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Monte Carlo Randomized Algorithm

■ Monte Carlo Randomized Algorithm (MCRA): A randomized algorithm that may produce incorrect results, but with bounded error probability

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Las Vegas Randomized Algorithm

■ Las Vegas Randomized Algorithm (LVRA): A randomized algorithm that always produces correct results, the only variation from one run to another is the running time

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Randomization of Uncertain Systems

- lacktriangle Consider random uncertainty Δ , associated pdf and bounding set B
- ∆ is a (real or complex) random vector (parametric uncertainty) or matrix (nonparametric uncertainty)
- Consider a performance function

$$J(\Delta): B \to \mathbf{R}$$

and level $\gamma > 0$

■ Define worst case and average performance

$$J_{\max} = \max_{\Delta \in B} J(\Delta)$$

 $J_{\text{ave}} = E_{\Lambda}(J(\Delta))$

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Example: H_{∞} Performance - 1

■ H_{∞} performance of sensitivity function

$$B = \{\Delta : \Delta = \text{bdiag } (\Delta_1, \dots, \Delta_q) \in \mathbf{F}^{n,m}, \, \sigma_{max}(\Delta) \leq \rho\}$$

$$S(s,\Delta) = 1/(1 + P(s,\Delta) C(s))$$

$$J(\Delta) = ||S(s,\Delta)||_{\infty}$$

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Example: H_{∞} Performance - 2

■ H_{∞} performance of sensitivity function

$$\begin{split} B &= \{\Delta \colon \Delta = \text{bdiag } (\Delta_1, \dots, \Delta_q) \in \mathbb{F}^{n,m}, \, \sigma_{\max}(\Delta) \leq \rho \} \\ S(s, \Delta) &= 1/(1 + P(s, \Delta) \, C(s)) \\ J(\Delta) &= \|S(s, \Delta)\|_{\infty} \end{split}$$

■ Objective: Check if

$$J_{\max} \leqslant \gamma$$
 and $J_{\text{ave}} \leqslant \gamma$

■ These are uncertain decision problems

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Two Problem Instances

■ We have two problem instances for worst case performance

$$J_{\mathrm{max}} \leqslant \gamma$$
 and $J_{\mathrm{max}} > \gamma$

and two problem instances for average case performance

$$J_{\mathrm{ave}} \leqslant \gamma \quad \mathrm{and} \quad J_{\mathrm{ave}} > \gamma$$

■ This leads to one-sided and two-sided MCRA

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One-Sided MCRA

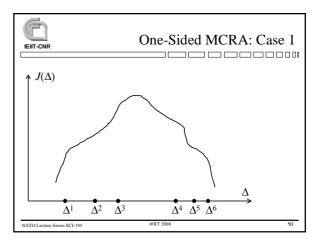
- One-sided MCRA: Always provide a correct solution in one of the instances (they may provide a wrong solution in the other instance)
- Consider the empirical maximum

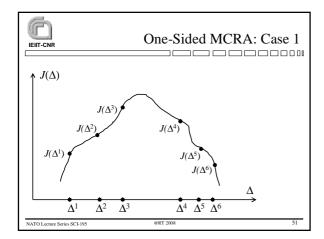
$$\hat{J}_{\max} = \max_{i=1,\dots,N} J(\Delta^i)$$

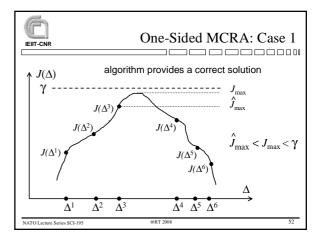
where Δ^i are random samples and N is the sample size

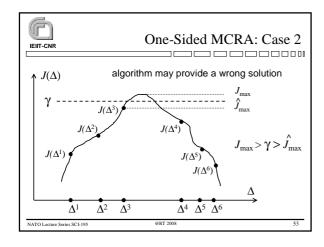
■ Check if $\hat{J}_{max} \leq \gamma$ or $\hat{J}_{max} > \gamma$

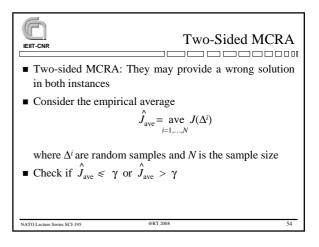
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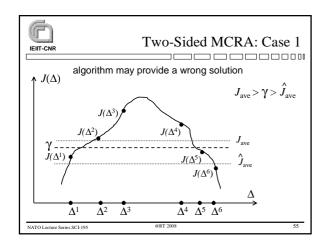


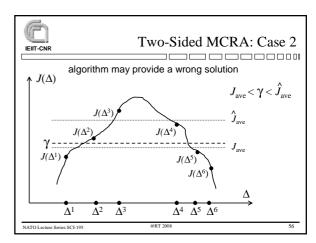










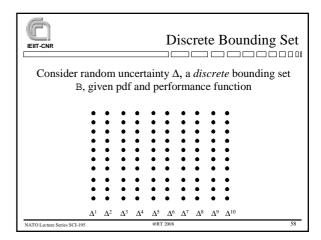




Las Vegas Randomized Algorithms

- We also have zero-sided (Las Vegas) randomized algorithms
- Las Vegas Randomized Algorithm (LVRA): Always give the correct solution
- The solution obtained with a LVRA is probabilistic, so "always" means with probability one
- Running time may be different from one run to another
- We can study the average running time

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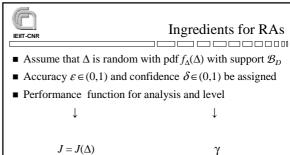




The Las Vegas Viewpoint

- Consider discrete random variables
- lacktriangle The sample space is discrete and M^N possible choices can be made
- In the binary case we have 2^N
- Finding maximum requires ordering the 2^N choices
- Las Vegas can be used for ordering real numbers
- Example: Randomized Quick Sort for sorting real numbers (classical in computer science)

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Randomized Algorithms for Analysis ______

- Two classes of randomized algorithms for probabilistic robust performance analysis
- P1: Performance verification (compute p_{γ})
- P2: Worst-case performance (compute J_{max})
- Both are based on uncertainty randomization of Δ
- Bounds on the sample size are obtained



Randomized Algorithms - 2

- \blacksquare We estimate p_{γ} by means of a randomized algorithm
- First, we generate N i.i.d. samples

$$\Delta^1, \Delta^2, ..., \Delta^N \in \mathcal{B}_D$$

according to the density f_{Δ}

■ We evaluate $J(\Delta^1), J(\Delta^2), ..., J(\Delta^N)$



Empirical Probability

■ Construct an indicator function

$$I(\Delta^i) = \begin{cases} 1 & \text{if } J(\Delta^i) \le \gamma \\ 0 & \text{otherwise} \end{cases}$$
• An estimate of p_{γ} is the empirical probability

$$\hat{p}_N = \frac{1}{N} \sum_{i=1}^{N} I(\Delta^i) = \frac{N_{good}}{N}$$

where N_{good} is the number of samples such that $J(\Delta^i) \leq \gamma$



A Reliable Estimate ____

■ The empirical probability is a reliable estimate if

$$|p_{\gamma} - \hat{p}_{N}| = |\Pr\{J(\Delta) \le \gamma\} - \hat{p}_{N}| \le \varepsilon$$

■ Find the minimum *N* such that

$$\Pr\{|p_{\gamma} - \hat{p}_{N}| \leq \varepsilon\} \geq 1 - \delta$$

where $\varepsilon \in (0,1)$ and $\delta \in (0,1)$



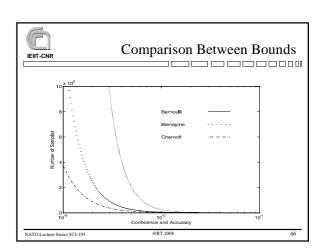
Chernoff Bound^[1]

■ For any $\varepsilon \in (0,1)$ and $\delta \in (0,1)$, if

$$N \ge \frac{\log \frac{2}{\delta}}{2\varepsilon^2}$$

then

$$\Pr\{|p_{\gamma} - \hat{p}_{N}| \leq \varepsilon\} \geq 1 - \delta$$





Chernoff Bound

- Remark: Chernoff bound improves upon other bounds such as Bernoulli (Law of Large Numbers)
- Dependence on $1/\delta$ is logarithmic
- Dependence on $1/\varepsilon$ is quadratic

ε	0.1%	0.1%	0.5%	0.5%
$1-\delta$	99.9%	99.5%	99.9%	99.5%
N	3.9.106	3.0.106	1.6.106	1.2.105

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Computational Complexity of RAs

- RAs are efficient (polynomial-time) because
- 1. Random sample generation of Δ^i can be performed in polynomial-time
- 2. Cost associated with the evaluation of $J(\Delta^i)$ for fixed Δ^i is polynomial-time
- 3. Sample size is polynomial in the problem size and probabilistic levels arepsilon and δ

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1. Random Sample Generation

- Random number generation (RNG): Linear and nonlinear methods for uniform generation in [0,1) such as Fibonacci, feedback shift register, BBS, MT, ...
- Non-uniform univariate random variables: Suitable functional transformations (e.g., the inversion method)
- The problem is much harder: Multivariate generation of samples of Δ with pdf $f_{\Delta}(\Delta)$ and support \mathcal{B}_D
- It can be resolved in polynomial-time

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2. Cost of Checking Stability

■ Consider a polynomial

$$p(s,a) = a_0 + a_1 s + \dots + a_n s^n$$

■ To check left half plane stability we can use the Routh test. The number of multiplications needed is

$$\frac{n^2}{4}$$
 for *n* even

$$\frac{n^2-1}{4}$$
 for n odd

- The number of divisions and additions is equal to this number
- We conclude that checking stability is $O(n^2)$

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3. Bounds on the Sample Size

- Chernoff bound is independent on the size of \mathcal{B}_D , on the structure D on the number of blocks, on the pdf $f_{\Lambda}(\Delta)$
- It depends only on δ and ϵ
- Same comments can be made for other bounds such as Bernoulli

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Worst-Case Performance

■ Recall that

$$J_{\max} = \max_{\Delta \in \mathcal{B}_D} J(\Delta)$$

■ Generate N i.i.d. samples

$$\Delta^1, \Delta^2, ..., \Delta^N \in \mathcal{B}_D$$

according to the density f_{Δ}

■ Compute the empirical maximum

$$\hat{J}_{\max} = \max_{i=1,\dots,N} J(\Delta^i)$$

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Worst-Case Bound (Log-over-Log)^[1]

■ For any $\varepsilon \in (0,1)$ and $\delta \in (0,1)$, if

$$N \ge \frac{\log \frac{1}{\delta}}{\log \frac{1}{1-\varepsilon}}$$

then

$$\Pr\{\Pr\{J(\Delta) > \hat{J}_N\} \le \varepsilon\} \ge 1 - \delta$$

[1] R. Tempo, E. W. Bai and F. Dabbene (1996)

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Comparison and Comments

- Number of samples is much smaller than Chernoff
- Bound is a specific instance of the fpras (fully polynomial randomized approximated scheme) theory
- Dependence on $1/\varepsilon$ is basically linear $\left(\log \frac{1}{1-\varepsilon} \approx \varepsilon\right)$

ε	0.1%	0.1%	0.5%	0.5%	0.01%	0.001%
1-δ	99.9%	99.5%	99.9%	99.5%	99.99%	99.999%
N	$6.91 \cdot 10^{3}$	5.30·10 ³	$1.38 \cdot 10^3$	$1.06 \cdot 10^3$	$9.21 \cdot 10^4$	1.16.106



Volumetric Interpretation

_____ ■ In the case of $f_{\Delta}(\Delta)$ uniform, we have

$$\Pr\{J(\Delta) > \hat{J}_N\} = \frac{vol(\Delta_{bad})}{vol(\mathcal{B}_D)}$$

■ Therefore

$$\Pr{\Pr{J(\Delta) > \hat{J}_N} \le \varepsilon} \ge 1 - \delta$$

is equivalent to

$$\Pr\{vol(\boldsymbol{\Delta}_{bad}) \leq \varepsilon vol(\mathcal{B}_D)\} \geq 1 - \delta$$



Confidence Intervals

■ The Chernoff and worst-case bounds can be computed *a*priori and provide an explicit functional relation

$$N = N(\varepsilon, \delta)$$

- The sample size obtained with the confidence intervals is not explicit
- Given $\delta \in (0,1)$, upper and lower confidence intervals p_L and p_U are such that $\Pr\{p_L \le p_\gamma \le p_U\} = 1 - \delta$



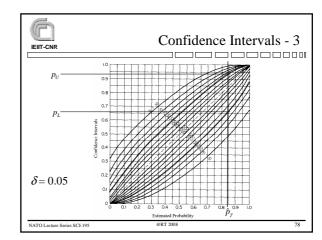
Confidence Intervals - 2 _____

■ The probabilities p_L and p_U can be computed a posteriori when the value of N_{good} is known, solving equations of the type

$$\begin{split} \sum_{k=N_{good}}^{N} \binom{N}{k} p_L^k (1-p_L)^{N-k} &= \delta_L \\ \sum_{k=0}^{N_{good}} \binom{N}{k} p_U^k (1-p_U)^{N-k} &= \delta_U \end{split}$$

$$\sum_{k=0}^{N_{good}} \binom{N}{k} p_U^k \left(1 - p_U\right)^{N-k} = \delta_U$$

with $\delta_L + \delta_U = \delta$





Statistical Learning Theory ______

■ The Chernoff Bound studies the problem $\Pr\{|p_{\gamma} - \hat{p}_{N}| \le \varepsilon\} \ge 1 - \delta$

where $p_{\gamma} = \Pr\{J(\Delta) \le \gamma \}$

- \blacksquare Performance function J is fixed
- Statistical Learning Theory computes bounds on the sample size for the problem

 $\Pr\{|\Pr(J(\Delta) \leq \gamma) - \hat{p}_N| \leq \varepsilon, \forall J \in \mathcal{J}\} \geq 1 - \delta$ where $\boldsymbol{\mathcal{J}}$ is a given class of functions

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VC and P-dimension^[1,2]

- Statistical Learning Theory aims at studying uniform Law of Large Numbers
- The bounds obtained depend on quantities called VCdimension (if J is a binary valued function), or Pdimension (if J is a continuous valued function)
- VC and P-dimension are measures of the problem complexity

[1] M. Vidyasagar (1997)

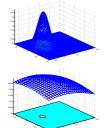
[2] E.D. Sontag (1998)

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Choice of the Distribution - 1

- The probability $Pr\{\Delta \in \mathcal{S}\}$ depends on $f_{\Lambda}(\Delta)$
- It may vary between 0 and 1 depending on the $\operatorname{pdf} f_{\Lambda}(\Delta)$





Choice of the Distribution - 2

- The bounds discussed are independent on the choice of the distribution but for computing $\Pr\{J(\Delta) \le \gamma\}$ we need to know the distribution $f_{\Lambda}(\Delta)$
- Some research has been done in order to find the worst-case distribution in a certain class^[1]
- Uniform distribution is the worst-case if a certain target is convex and centrally symmetric

[1] B. R. Barmish and C. M. Lagoa (1997)

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Choice of the Distribution - 3 _____

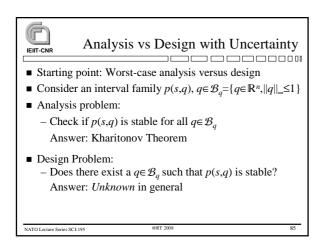
■ Minimax properties of the uniform distribution have been studied[1]

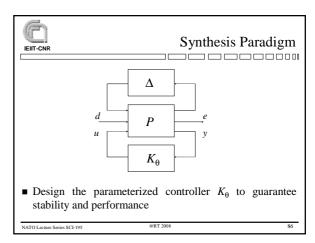
[1] E. W. Bai, R. Tempo and M. Fu (1998)

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Probabilistic Robust Synthesis







Synthesis Performance Function

- Recall that the parameterized controller is K_{Θ}
- We replace $J(\Delta)$ with a synthesis performance function

$$J = J(\Delta, \theta)$$

where $\theta \in \Theta$ represents the controller parameters to be determined and their bounding set

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Randomized Algorithms for Synthesis

- Two classes of RAs for probabilistic synthesis
- Average performance synthesis^[1]
- Based on expected value minimization
- Use of Statistical Learning Theory results
- Very general problems can be handled
- Existing bounds are very conservative and controller randomization is required
- Ongoing research aiming at major reduction of sample size

[1] M. Vidyasagar (1998

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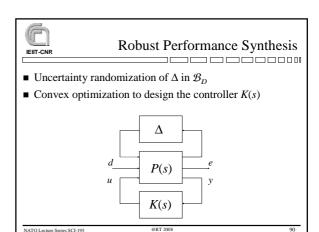


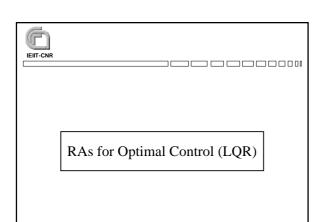
Randomized Algorithms for Synthesis

- Robust performance synthesis^[1]
- Problem reformulation as robust feasibility
- Only convex problems can be handled
- Finite-time convergence with probability one is obtained

[1]B. Polyak and R. Tempo (2001)

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Uncertain Systems in State Space

■ We consider a state space description of the uncertain system

$$\dot{x}(t) = A(\Delta)x(t) + Bu(t)$$

with $x(0)=x_0$; $x \in \mathbb{R}^n$; $u \in \mathbb{R}^m$, $\Delta \in \mathcal{B}_D$

■ For example, $A(\Delta)$ is an interval matrix with bounded entries $a_{ii}^- \le a_{ii} \le a_{ii}^+$

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Interval and Vertex Matrices

- That is, a_{ik} ranges in the interval for all i, k

$$|a_{ik} - a_{ik}^*| \leq w_{ik}$$

■ We consider interval uncertainty A (i.e. when $\Delta \in \mathcal{B}_D$)

where a_{ik}^{*} are nominal values and w_{ik} are weights

■ Define the $N = 2^{n^2}$ vertex matrices $A^1, A^2, ..., A^N$

$$a_{ik}={a_{ik}}^*+w_{ik}\quad\text{or}\quad a_{ik}={a_{ik}}^*-w_{ik}$$
 for all $i,\,k=1,\,2,\,\ldots,\,n$

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Common Lyapunov Functions

■ Given matrices *A**, *W* and feedback *K*, find a *common* quadratic Lyapunov function *Q* > 0 for the system

$$\dot{x}(t) = (A + B K) x(t)$$
 for all $A \in A$

■ Find Q > 0 such that

$$L(Q, A) = (A+BK)^T Q + Q (A+BK) < 0$$
 for all $A \in A$

■ Equivalently, find Q > 0 such that

$$\lambda_{max} L(Q, A) < 0$$
 for all $A \in A$

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Lyapunov Stability of Interval Systems

- Quadratic Lyapunov stability analysis and synthesis of interval systems are NP-hard problems
- In principle, they can be solved in one-shot with convex optimization, but the number of constraints is exponential
- We can use relaxation (e.g. $\pi/2$ Theorem^[1]) or randomization

[1] Yu. Nesterov (1997)

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Vertex Solution

- Due to convexity, it suffices to study *L(Q, A)* < 0 for all vertex matrices^[1]
- Question: Do we really need to check all the vertex matrices $(N = 2^{n^2})$?

[1] H.P. Horisberger, P.R. Belanger (1976)

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Vertex Reduction

- Answer: It suffices to check "only" a subset of 2²ⁿ vertex matrices^[1]
- This is still exponential (the problem is NP-hard), but it leads to a major computational improvement for medium size problems (e.g. *n* = 8 or 10)
- For example, for n=8, N is of the order 10^5 (instead of 10^{19})

[1] T. Alamo, R. Tempo, D. Rodriguez, E.F. Camacho (2007)

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Diagonal Matrices and Generalizations

- Transform the original problem from full square matrices A to diagonal matrices $Z \in \mathbb{R}^{2n,2n}$
- It suffices to check the vertices of Z
- lacktriangle Extensions for L_2 -gain minimization and other related LMI problems
- Generalizations for multiaffine interval systems

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Las Vegas Randomized Algorithm

- We may perform randomization of the $N = 2^{n^2}$ vertices (in the worst case)
- If we select the vertices in random order according to a given pdf, we have a LVRA

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Probabilistic Solution

■ Randomly generate $A^1,...,A^N$. Then, check if the Lyapunov equation

$$A^iQ + Q(A^i)^T \le 0$$

is feasible for i=1,...,N and find a common solution $Q=Q^T>0$

lacktriangle Critical problem: Even if N is relatively small, this is a hard computational problem

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Sequential Algorithm

- Key point: Sequential algorithm which deals with one constraint at each step
- \blacksquare At step k we have

Phase 1: Uncertainty randomization of Δ

Phase 2: Gradient algorithm and projection

■ Final result: Find a solution $Q=Q^T>0$ with probability one in a finite number of steps

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Definition

■ Let \mathcal{E}_n be an Euclidean space

$$\mathcal{E}_n = \left\{ A = A^T \in \mathbb{R}^n, ||A|| = \sqrt{\sum_{i,k=1}^n a_{1k}^2} \right\}$$

and C be the cone of positive semi-definite matrices

$$C = \{ A \in \mathcal{E}_n : A \ge 0 \}$$

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Projection on a Cone

■ For any real symmetric matrix A we define the projection $[A]^+ \in C$ as

$$[A]^+ = \arg\min_{X \in C} ||A - X||$$

- The projection can be computed through the eigenvalue decomposition $A=T\Lambda T^T$
- Then

$$[A]^{+} = T\Lambda^{+}T^{T}$$

where $\lambda_i^+ = \max \{\lambda_i, 0\}$

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Phase 1: Uncertainty Randomization

- Uncertainty randomization: Generate $\Delta^k \in \mathcal{B}_D$
- Then, for guaranteed cost we obtain the Lyapunov equation

$$A(\Delta^k)Q + QA^T(\Delta^k) \leq 0$$

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Matrix Valued Function

■ Define a matrix valued function

$$V(Q, \Delta^k) = A(\Delta^k)Q + QA^T(\Delta^k)$$

and a scalar function

$$v(Q, \Delta^k) = \| [V(Q, \Delta^k)]^{\dagger} \|$$

where $\|\cdot\|$ is the Frobenius norm

■ We can also take the maximum eigenvalue of $V(Q, \Delta^k)$

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Phase 2: Gradient Algorithm

■ We write

$$Q^{k+1} = \begin{cases} \left[Q^k - \mu^k \partial_Q \left\{ \nu \left(Q^k, \Delta^k \right) \right\} \right]^{\dagger} & \text{if } \nu \left(Q^k, \Delta^k \right) > 0 \\ Q^k & \text{otherwise} \end{cases}$$

where ∂_Q is the subgradient and the stepsize μ^k is

$$\mu^{k} = \frac{v\left(Q^{k}, \Delta^{k}\right) + r\left\|\partial_{Q}\left\{v\left(Q^{k}, \Delta^{k}\right)\right\}\right\|}{\left\|\partial_{Q}\left\{v\left(Q^{k}, \Delta^{k}\right)\right\}\right\|^{2}}$$

and r>0 is a parameter

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Closed-form Gradient Computation

■ The function $\nu(Q, \Delta^k)$ is convex in Q and its subgradient can be easily computed in a closed form

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Theorem^[1]

- Assumption: Every open subset of \mathcal{B}_D has positive
- Theorem: A solution Q, if it exists, is found in a finite number of steps with probability one
- Idea of proof: The distance of Q^k from the solution set decreases at each correction step

[1] B.T. Polyak and R. Tempo (2001)

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Example^[1]

- We study a multivariable example for the design of a controller for the lateral motion of an aircraft.
- The model consists of four states and two inputs

$$\dot{x}(t) = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & L_{\rho} & L_{\beta} & L_{r} \\ \frac{g}{\sqrt{r}} & 0 & Y_{\beta} & -1 \\ N_{\beta}(\frac{g}{\sqrt{r}}) & N_{\rho} & N_{\beta} + N_{\beta}Y_{\beta} & N_{r} - N_{\beta} \end{bmatrix} x(t) + \begin{bmatrix} 0 & 0 \\ 0 & -3.91 \\ 0.035 & 0 \\ -2.53 & 0.31 \end{bmatrix} u(t)$$

[1] B.D.O. Anderson and J.B. Moore (1971)

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Example - 2 ____

- The state variables are
 - $-x_1$ bank angle
 - $-x_2$ derivative of bank angle
 - $-x_3$ sideslip angle
 - x_4 jaw rate
- The control inputs are
 - $-u_1$ rudder deflection
 - u2 aileron deflection



Example - 3

- ____ values: $L_p = -2.93$, $L_{\beta} = -4.75$, $L_r = 0.78$, Nominal g/V=0.086, $Y_{\beta}=-0.11$, $N_{\beta}=0.1$, $N_{\rho}=-0.042$, $N_{\beta}=2.601$, $N_r = -0.29$
- Perturbed matrix $A(\Delta)$: each parameter can take values in a range of ±15% of the nominal value
- Quadratic stability (γ =0): take R=I and S=0.01I
- Remark: $A(\Delta)$ is multiaffine in the uncertain parameters: quadratic stability can be ascertained solving simultaneously 29=512 LMIs



Example - 4 ----

- Sequential algorithm:
 - Initial point Q_0 randomly selected
 - -800 random matrices Δ^k
 - The algorithm converged to

$$Q = \begin{bmatrix} 0.7560 & -0.0843 & 0.1645 & 0.7338 \\ -0.0843 & 1.0927 & 0.7020 & 0.4452 \\ 0.1645 & 0.7020 & 0.7798 & 0.7382 \\ 0.7338 & 0.4452 & 0.7382 & 1.2162 \end{bmatrix}$$



Example - 5

- The corresponding controller
 - $K = B^T Q^{-1} = \begin{bmatrix} 38.6191 & -4.3731 & 43.1284 & -49.9587 \end{bmatrix}$ -2.8814 -10.1758 10.2370 -0.4954 satisfies all the 512 vertex LMIs and therefore it is also a quadratic stabilizing controller in a deterministic sense
- The optimal LQ controller computed on the nominal plant satisfies only 240 vertex LMIs

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Extensions



Related Literature and Extensions

- Minimization of a measure of violation for problems that are not strictly feasible^[1]
- Uncertainty in the control matrix, $B=B(\Delta)$, $\Delta \in \mathcal{B}_D$ We take the feedback law

$$u = YO^{-1}x$$

where Y and $Q=Q^T>0$ are design variables

[1] B.R. Barmish and P. Shcherbakov (1999)

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Related Literature

- Related literature on optimization and adaptive control with linear constraints^[1,2,3,4]
- Stochastic approximation algorithms have been widely studied in the stochastic control and optimization literature[6,7]

[1] S. Agmon (1954)

[2] T.S. Motzkin and I.J. Schoenberg (1954)[3] B.T. Polyak (1964)

[4] V.A. Bondarko and V.A. Yakubovich (1992)

[6] H.J. Kushner and G.G. Yin (2003)

[7] J.C. Spall (2003)



Subsequent Research

- _____ ■ Design of common Lyapunov functions for switched systems[1]
- From common to piecewise Lyapunov functions^[2]
- Ellipsoidal algorithm instead of gradient algorithm^[3]
- Stopping rule which provides the number of steps^[4]
- Other algorithms have been recently proposed^[5-6]
- [1] D. Liberzon and R. Tempo (2004) [2] H. Ishii, T. Basar and R. Tempo (2005) [3] S. Kanev, B. De Schutter and M. Verhaegen (2002)

- [4] Y. Oishi and H. Kimura (2003) [5] Y. Fujisaki and Y. Oishi (2007) [6] T. Alamo, R. Tempo, D. R. Ramirez and E. F. Camacho (2007)

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Optimization Problems^[1]

- ■ Extensions to optimization problems
- Consider convex function f(x) and function $g(x,\Delta)$ convex in x for fixed Δ
- Semi-infinite (nonlinear) programming problem

 $\min f(x)$

 $g(x,\Delta) \le 0$ for all $\Delta \in \mathcal{B}$

- Reformulation as stochastic optimization
- Drawback: Convergence results are only asymptotic

[1] V. B. Tadic, S. P. Meyn and R. Tempo (2003)



Scenario Approach

- _____i ■ The scenario approach for convex problems^[1]
- Non-sequential method which provides a one-shot solution for general convex problems
- Randomization of $\Delta \in \mathcal{B}$ and solution of a single convex optimization problem
- Derivation of a bound on the sample size^[1]
- A new improved bound based on a pack-based strategy^[2]

[1] G. Calafiore and M. Campi (2004)

[2] T. Alamo, R. Tempo and E.F. Camacho (2007)

Convex Semi-Infinite Optimization

■ The semi-infinite optimization problem is

 $\min c^T \theta$ subject to $f(\theta, \Delta) \le 0$ for all $\Delta \in \mathcal{B}$

where $f(\theta, \Delta) \le 0$ is convex in θ for all $\Delta \in \mathcal{B}$

- We assume that this problem is either unfeasible or, if feasible, it attains a unique solution for all $\Delta \in \mathcal{B}$ (this assumption is technical and may be removed)
- We assume that $\theta \in \Theta \subseteq \mathbb{R}^n$



Scenario Problem

- _____ ■ Using randomization, we construct a scenario problem
- Taking random samples Δ^i , i = 1, 2, ..., N, we construct

 $f(\theta, \Delta^i) \le 0, \quad i = 1, 2, ..., N$

and

min $c^T \theta$ subject to $f(\theta, \Delta^i) \le 0$, i = 1, 2, ..., N

Theorem^[1]

■ Theorem: For any $\varepsilon \in (0,1)$ and $\delta \in (0,1)$, if

 $N \ge \left[\frac{2}{\varepsilon} \log(1/\delta) + 2n + \frac{2n}{\varepsilon} \log(2/\varepsilon) \right]$ then, with probability no smaller than 1- δ

- either the scenario problem is unfeasible and then also the semi-infinite optimization problem is unfeasible
- or, the scenario problem is feasible, then its optimal solution $\hat{\theta}_N$ satisfies

$$\Pr\{ \Delta \in \mathcal{B} : f(\theta, \Delta) > 0 \} \leq \mathcal{E}$$

[1] G. Calafiore and M. Campi (2004)



A New Improved Bound^[1]

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■ A new improved bound (based on a so-called packbased strategy) has been recently obtained

 $N \ge \left[2/\varepsilon \log(1/2\delta) + 2n + 2n/\varepsilon \log 4 \right]$

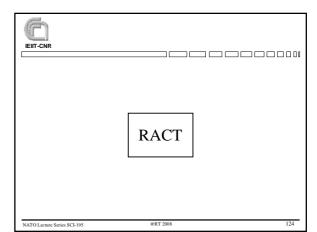
■ The main difference with the previous bound is that the

 $2n/\varepsilon \log (2/\varepsilon)$

is replaced with

 $2n/\varepsilon \log 4$

[1] T. Alamo, R. Tempo and E.F. Camacho (2007)





RACT

- RACT: Randomized Algorithms Control Toolbox for Matlab
- RACT has been developed at IEIIT-CNR and at the Institute for Control Sciences-RAS, based on a bilateral international project
- Members of the project

Andrey Tremba (Main Developer and Maintainer)

Giuseppe Calafiore

Fabrizio Dabbene

Elena Gryazina Boris Polyak (Co-Principal Investigator)

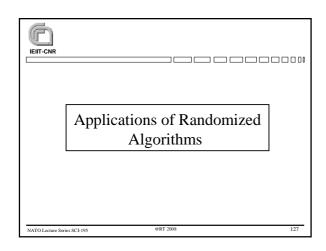
Pavel Shcherbakov

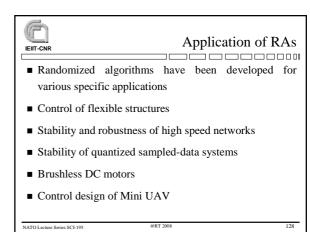
Roberto Tempo (Co-Principal Investigator)

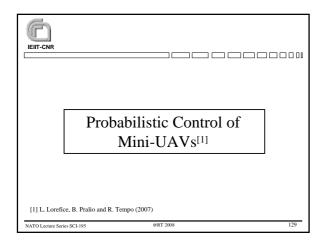
RACT

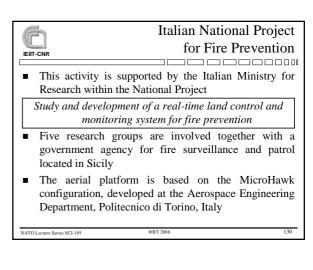
- Main features
- Define a variety of uncertain objects: scalar, vector and matrix uncertainties, with different pdfs
- Easy and fast sampling of uncertain objects of almost any type
- Randomized algorithms for probabilistic performance verification and probabilistic worst-case performance
- Randomized algorithms for feasibility of uncertain LMIs using stochastic gradient, ellipsoid or cutting plane methods (YALMIP needed)

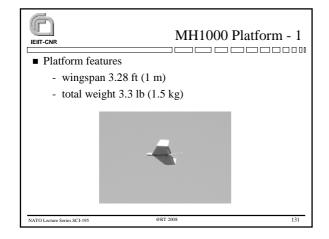
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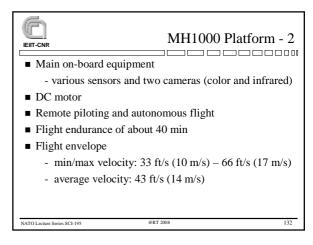


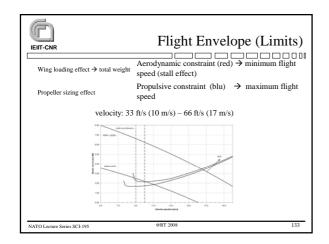


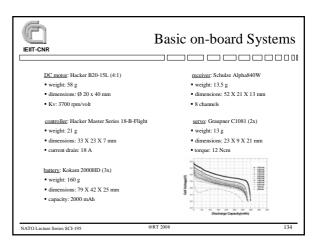


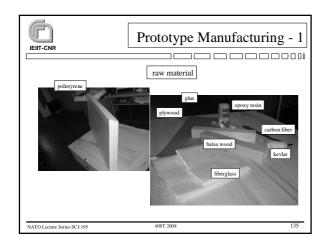


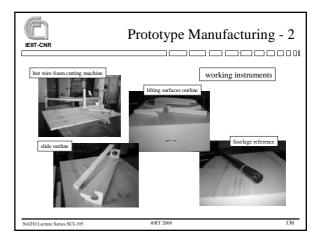


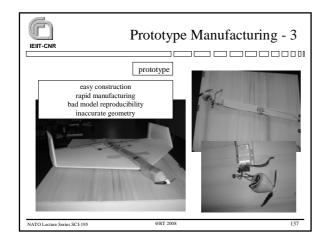


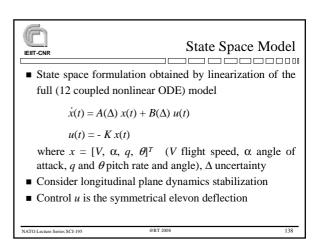














Uncertainty Description - 1

- We consider structured parameter uncertainties affecting plant and flight conditions, and aerodynamic database
- Uncertainty vector $\Delta = [\delta_1, ..., \delta_{16}]$ where $\delta_i \in [\delta_i^-, \delta_i^+]$
- Key point: There is no explicit relation between state space matrices A and B and uncertainty Δ
- This is due to the fact that state space system is obtained through linearization and off-line flight simulator
- The only techniques which could be used in this case are simulation-based which lead to randomized algorithms

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Uncertainty Description - 2

- We consider random uncertainty $\Delta = [\delta_1, ..., \delta_{16}]^T$
- The pdf is either uniform (for plant and flight conditions) or Gaussian (for aerodynamic database uncertainties)
- Flight conditions uncertainties need to take into account large variations on physical parameters
- Uncertainties for aerodynamic data are related to experimental measurement or round-off errors

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Plant and Flight Condition Uncertainties

parameter	pdf	$\overline{\delta_i}$	%	δ_i^{-}	$\delta_i^{\scriptscriptstyle +}$	#
flight speed [ft/s]	U	42.65	± 15	36.25	49.05	1
altitude [ft]	U	164.04	± 100	0	328.08	2
mass [lb]	U	3.31	± 10	2.98	3.64	3
wingspan [ft]	U	3.28	± 5	3.12	3.44	4
mean aero chord [ft]	U	1.75	± 5	1.67	1.85	5
wing surface [ft ²]	U	5.61	± 10	5.06	6.18	6
moment of inertia [lb ft2]	U	1.34	± 10	1.21	1.48	7

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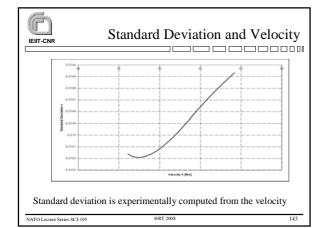


Aerodynamic Database Uncertainties

parameter	pdf	$\overline{\delta_{i}}$	σ_{i}	#
C_X [-]	G	-0.01215	0.00040	8
C _z [-]	G	-0.30651	0.00500	9
C _m [-]	G	-0.02401	0.00040	10
C_{Xq} [rad-1]	G	-0.20435	0.00650	11
C_{Zq} [rad-1]	G	-1.49462	0.05000	12
C_{mq} [rad-1]	G	-0.76882	0.01000	13
C_X [rad-1]	G	-0.17072	0.00540	14
C_Z [rad-1]	G	-1.41136	0.02200	15
C_m [rad-1]	G	-0.94853	0.01500	16

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Critical Parameters and Matrices

- We select flight speed (δ_1) and take off mass (δ_3) as critical parameters
- Flight speed is taken as critical parameter to optimize gain scheduling issues
- Take off mass is a key parameter in mission profile definition
- We define critical matrices

 $A_c^{\ 1}$ $A_c^{\ 2}$ $A_c^{\ 3}$ $A_c^{\ 4}$ $B_c^{\ 1}$ $B_c^{\ 2}$ $B_c^{\ 3}$ $B_c^{\ 4}$

■ They are constructed setting δ_1 , δ_3 to the extreme values δ_1 , δ_1^+ , δ_1^+ , δ_3^+ and all the remaining δ_i are equal to $\overline{\delta_i}$

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Phase 1: Random Gain Synthesis (RGS)

- ■ Critical parameters are flight speed and take off mass
- Specification property

$$S_1 = \{K: A_c - B_c K \text{ satisfies the specs below}\}$$

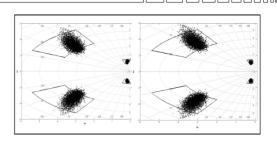
 $\omega_{SP} \in [4.0, 6.0] \text{ rad/s}$ $\zeta_{SP} \in [0.5, 0.9]$ $\omega_{PH} \in [1.0, 1.5] \text{ rad/s}$ $\zeta_{PH}\in[0.1,\!0.3]$ $\Delta\omega_{PH} < \pm~20\%$ $\Delta\omega_{SP}$ < \pm 45%

where ω and ζ are undamped natural frequency and damping ratio of the characteristic modes; SP and PH denote short period and phugoid mode

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Specs in the Complex Plane





Volume of the Good Set

■ Define a bounding set *B* of gains *K* _____

$$B = \{K: k_i \in [k_i, k_i], i = 1,...,4\}$$

■ Define the volume of the *good set*

$$\operatorname{Vol}_{good} = \int_A dK$$

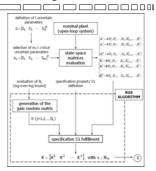
where $A = \{ K \in B \cap S_1 \}$

• Vol_B is simply the volume of the hyperrectangle B



Randomized Algorithm 1 (RGS)

- Uniform pdf for controller gains K in given intervals
- Accuracy and confidence ϵ =4 ·10⁻⁵ and δ = 3 · 10⁻⁴
- Number of random samples is computed with "Log-over-Log" Bound obtaining N = 200,000
- We obtained 5 gains K^i satisfying specification property S_1





Randomized Algorithm 1 (RGS)

Given ϵ , $\delta \in (0,1)$, RGS returns the set of gains $\{K^1, ..., K^s\}$ satisfying S_1

- Compute N using the Log-over-log Bound;
- For fixed j=1,2,...,N, generate uniformly the gain random matrix $K^{j} \in B$;
- Set *C*=0;
- For fixed i=1,2,3,4, compute the closed-loop matrix $A_{c^i}(K^j) = A_{c^i}^i B_{c^i}^i K^j;$ $\text{if } K^j \in S_1, \text{ set } C = C+1;$

- otherwise, set C = C;

- 5. End;
- If C = 4, return the gain K^{j} ;
- Set j = j+1 and return to Step 2;



Random Gain Set _____

gain set	K_V	K_{α}	K_q	K_{θ}
K^1	0.00044023	0.09465000	0.01577400	-0.00473510
K ²	0.00021450	0.09581200	0.01555500	-0.00323510
K ³	0.00054999	0.09430800	0.01548200	-0.00486340
K ⁴	0.00010855	0.09183200	0.01530000	-0.00404380
K ⁵	0.00039238	0.09482700	0.01609300	-0.00417340



Phase 2: Random Stability Robustness Analysis (RSRA) _____

- Take $K_{rand} = K^i$ obtained in Phase 1
- lacktriangledown Randomize Δ according to the given pdf and take Nrandom samples Δ^i
- Specification property

$$S_2 = \{\Delta: A(\Delta) - B(\Delta) \mid K_{rand} \text{ satisfies the specs of } S_1\}$$

■ Computation of the empirical probability of stability \hat{p}_N

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Empirical Probability

- Consider fixed gain K_{rand}
- Define the probability

$$p_{mn} = \int_C p(\Delta) d\Delta$$

 $p_{rue} = \int_{\mathcal{C}} p(\Delta) d \Delta$ where $\mathcal{C} = \{ \Delta \in B \cap S_2 \}$ and $p(\Delta)$ is the given pdf

Then, we introduce a "success" indicator function

$$I(\Delta i) = 1 \text{ if } \Delta i \in S_2$$

or $I(\Delta^{i}) = 0$ otherwise

■ The empirical probability for S_2 is given by

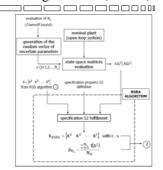
$$\hat{p}_N = N_{good}/N$$

where N_{good} is equal to the number of successes



Randomized Algorithm 2 (RSRA)

- Take K_{rand} from Phase 1
- Accuracy and confidence $\varepsilon = \delta = 0.0145$
- Number of random samples is computed with Chernoff Bound obtaining N = 5,000
- Empirical probability is defined using an indicator function



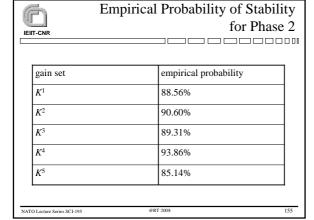
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Randomized Algorithm 2 (RSRA)

Given ε , $\delta \in (0,1)$, RSRA returns the empirical probability \hat{p}_N that S_2 is satisfied for a gain K_{rand} provided by Algorithm 1

- Compute N using the Chernoff Bound;
- 2. Generate N random vectors $\Delta \in B$ according to the given pdf;
- 3. For fixed j=1,2,...,N, compute the closed-loop matrix $A_{c,l}(\Delta^l) = A(\Delta^l) - B(\Delta^l)K_{rand},$ - if $A_{c,l}(\Delta^l) \in S_2$, set $I(\Delta^l) = 1$; - otherwise.

 - otherwise, set $I(\Delta \vec{v}) = 0$;
- 4. End:
- 5. Return the empirical probability \hat{p}_N





Probability Degradation Function

■ Flight condition uncertainties are multiplied by the amplification factor $\rho > 0$ keeping the nominal value constant

 $\delta_i \in \rho \left[\delta_i^-, \delta_i^+ \right] \quad \text{for } i = 1, 2, ..., 7$

■ No uncertainty affects the aerodynamic database, i.e.

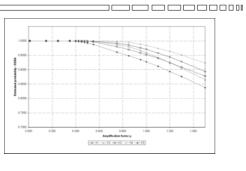
 $\delta_i = \overline{\delta_i}$ for i = 8, 9, ..., 16

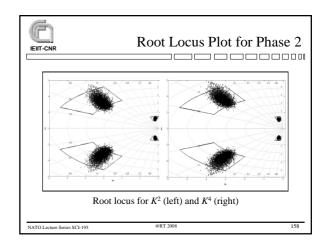
- For fixed $\rho \in [0,1.5]$ we compute the empirical probability for different gain sets K^i
- The plot empirical probability vs ρ is the probability degradation function

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Probability Degradation Function for Phase 2







Phase 3: Random Performance Robustness Analysis (RPRA)

- This phase is similar to Phase 2, but military specs are considered (bandwidth criterion)
- Specification property

 $S_3 = \{\Delta: A(\Delta) - B(\Delta) \mid K_{rand} \text{ satisfies the specs below}\}$

 $\omega_{BW} \in [2.5,5.0] \text{ rad/s}$ $\tau_P \in [0.0,0.5] \text{ s}$

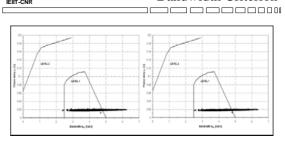
where ω_{BW} and τ_{P} are bandwidth and phase delay of the frequency response

lacksquare Computation of the empirical probability that S_3 is satisfied

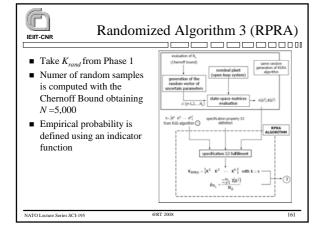
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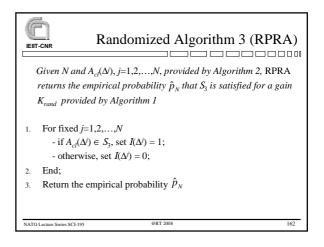


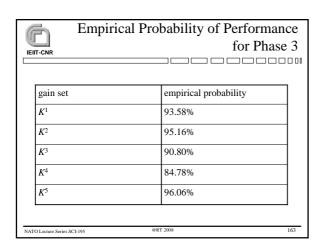
Bandwidth Criterion

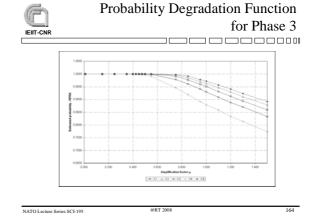


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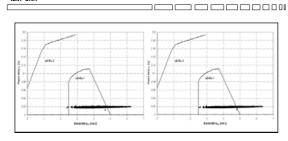






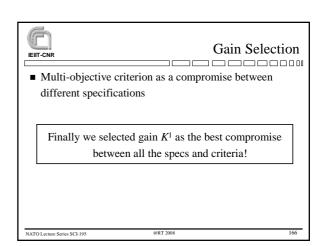


Bandwidth Criterion for Phase 3



Bandwidth criterion for K^1 (left) and K^3 (right)

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Conclusions: Flight Tests in Sicily - 1

- Evaluation of the payload carrying capabilities and autonomous flight performance
- Mission test involving altitude, velocity and heading changing was performed in Sicily
- Checking effectiveness of the control laws for longitudinal and lateral-directional dynamics
- Flight control design based on RAs for stabilization and guidance

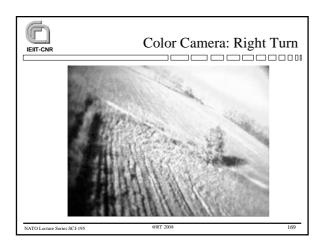
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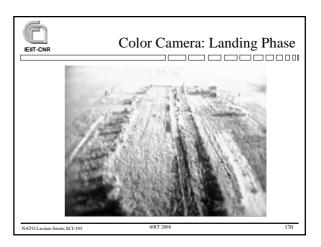


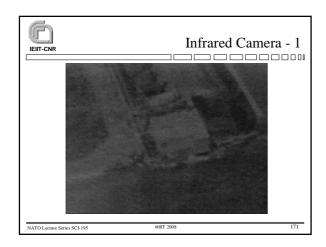
Conclusions: Flight Tests in Sicily - 2

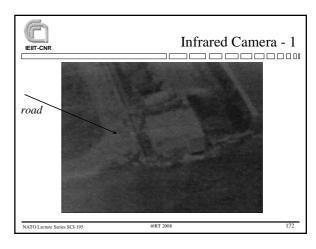
- Satisfactory response of MH1000
- Possible improvements by iterative design procedure
- Stability of the platform is crucial for the video quality and in the effectiveness of the surveillance and monitoring tasks

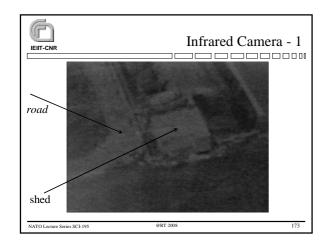
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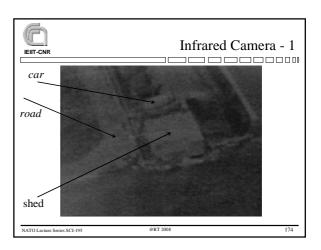


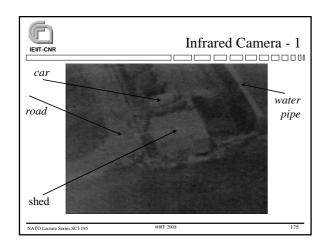


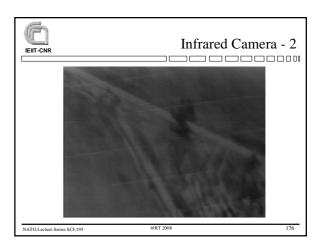


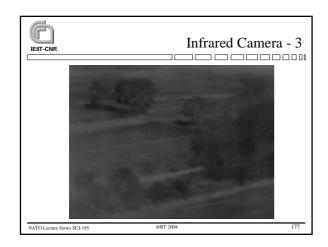


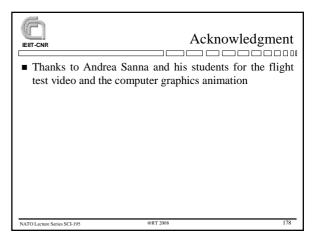


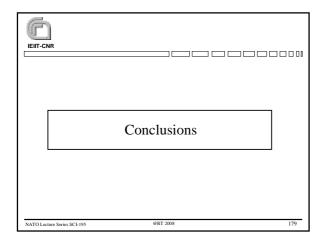


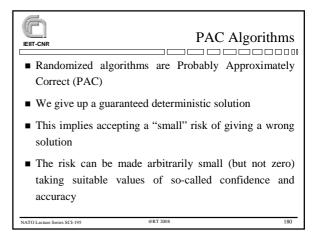












■ Two open problems	PAC Algorithms
 Optimization with sequential met Derive "reasonable" bounds for theory approach 	